

Database Management Systems

Triggers



Triggers

- \square Active Database Systems
- \sum Oracle Triggers
- \sum DB2 Triggers
- \sum Differences between Oracle and DB2
- \sum Trigger Design





Database Management Systems

Active Database Systems



Active database systems

 \sum Traditional DBMS operation is *passive*

- Queries and updates are explicitly requested by users
- The knowledge of processes operating on data is typically embedded into applications
- \sum *Active* database systems
 - Reactivity is a service provided by a normal DBMS
 - Reactivity *monitors* specific database events and *triggers* actions in response



Active database systems

 \sum Reactivity is provided by automatically executing rules

- \sum Rules are in the form
 - Event
 - Condition
 - Action
- $\mathop{\textstyle \sum}\nolimits$ Also called active or ECA rules



Active rules

\supset Event

- Database modification operation
- \supset Condition
 - Predicate on the database state
 - If the condition is true, the action is executed
- \supset Action
 - Sequence of SQL instructions or application procedure



Rule engine

 \sum Component of the DBMS, in charge of

- Tracking events
- Executing rules when appropriate
 - based on the execution strategy of the DBMS
- \sum Rule execution is interleaved with traditional transaction execution



Example

- \sum The active rule manages reorder in an inventory stock
 - when the stocked quantity of a product goes below a given threshold
 - a new order for the product should be issued
- ightarrow Event
 - Update of the quantity on hand for product x
 - Insert of a new product x



Example

- \sum The active rule manages reorder in an inventory stock
 - when the stocked quantity of a product goes below a given threshold
 - a new order for the product should be issued
- \supset Condition
 - The quantity on hand is below a given threshold *and* there are no pending orders for product x
- \supset Action
 - Issue an order with given reorder quantity for product x



Applications of active rules

\sum Internal applications

- maintenance of complex integrity constraints
- replication management
- materialized view maintenance
- \supset Business Rules
 - Incorporate into the DBMS application knowledge
 - E.g., reorder rule

 \supset Alerters

• widely used for notification



Triggers

- \sum Commercial products implement active rules by means of *triggers*
- \sum SQL provides instructions for defining triggers
 - Triggers are defined by means of the DDL instruction CREATE TRIGGER
- \sum Trigger syntax and semantics are covered in the SQL3 standard
 - Some commercial products implement different features with respect to the standard



Trigger structure

\supset Event

- Insert, delete, update of a table
- Each trigger can only monitor events on a *single* table
- \supset Condition
 - SQL predicate (it is optional)
- \supset Action
 - Sequence of SQL instructions
 - Proprietary programming language blocks
 - e.g. Oracle PL/SQL
 - Java block



Execution process

When the events take placeIf the condition is trueThen the action is executed

 \sum Seems very simple but...

- Execution modes
- Execution granularity

[triggering] [evaluation] [execution]



Execution mode

\supset Immediate

- The trigger is executed *immediately before* or *after* the triggering statement
- \supset Deferred

The trigger is executed immediately *before commit* Donly the immediate option is available in commercial systems



Execution granularity

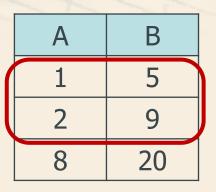
\sum Tuple (or row level)

- One separate execution of the trigger *for each tuple* affected by the triggering statement
- \supset Statement
 - One single trigger execution *for all tuples* affected by the triggering statement



Granularity example

Σ Table T



 \Box Transaction statement

UPDATE T SET A=A+1 WHERE B<10;

\sum Trigger execution

- A row level trigger executes twice
- A statement level trigger executes once



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Oracle Triggers



CREATE TRIGGER *TriggerName Mode Event* {OR *Event* } ON *TargetTable* [[REFERENCING *ReferenceName*] FOR EACH ROW [WHEN *Predicate*]] PL/SQL Block



CREATE TRIGGER TriggerName

Mode Event {OR *Event* }

ON TargetTable

[[REFERENCING *ReferenceName*] FOR EACH ROW [WHEN *Predicate*]] PL/SQL Block

\supset *Mode* is BEFORE or AFTER

Also INSTEAD OF but should be avoided



CREATE TRIGGER TriggerName

Mode Event {OR *Event* }

ON TargetTable

[[REFERENCING *ReferenceName*] FOR EACH ROW [WHEN *Predicate*]] PL/SQL Block

∑ Event **ON** *TargetTable* is

- INSERT
- DELETE
- UPDATE [OF ColumnName]



CREATE TRIGGER *TriggerName Mode Event* {OR *Event* } ON *TargetTable* [[REFERENCING *ReferenceName*] FOR EACH ROW [WHEN *Predicate*]] PL/SQL Block

 \sum FOR EACH ROW specifies row level execution semantics

• If omitted, the execution semantics is statement level



CREATE TRIGGER *TriggerName Mode Event* {OR *Event* } ON *TargetTable* [[REFERENCING *ReferenceName*] FOR EACH ROW [WHEN *Predicate*]] PL/SQL Block

 \sum The old and new states of the row triggering a *row level* trigger may be accessed by means of the

- OLD. ColumnName variable
- NEW.ColumnName variable



CREATE TRIGGER *TriggerName Mode Event* {OR *Event* } ON *TargetTable* [[REFERENCING *ReferenceName*] FOR EACH ROW [WHEN *Predicate*]] PL/SQL Block

To rename the state variables
 REFERENCING OLD AS *OldVariableName* similarly for NEW



CREATE TRIGGER *TriggerName Mode Event* {OR *Event* } ON *TargetTable* [[REFERENCING *ReferenceName*] FOR EACH ROW [WHEN *Predicate*]]

PL/SQL Block

 \sum *Only* for row level execution semantics (i.e., FOR EACH ROW)

- A condition may be optionally specified
- The old and new state variables may be accessed



CREATE TRIGGER *TriggerName Mode Event* {OR *Event* } ON *TargetTable* [[REFERENCING *ReferenceName*] FOR EACH ROW [WHEN *Predicate*]] PL/SQL Block

 \supset The action is

- a sequence of SQL instructions
- a PL/SQL block

 \sum No transactional and DDL instructions



Trigger semantics

- \supset Execution modes
 - immediate before
 - immediate after
- Σ Granularity is
 - row (tuple)
 - statement
- \sum Execution is triggered by insert, delete, or update statements in a transaction



Execution algorithm

- 1. Before statement triggers are executed
- 2. For each tuple in *TargetTable* affected by the triggering statement
 - a) Before row triggers are executed
 - b) The triggering statement is executed+ integrity constraints are checked on tuples
 - c) After row triggers are executed
- 3. Integrity constraints on tables are checked
- 4. After statement triggers are executed



Trigger semantics

 ${\hfill}{>}$ The execution order for triggers with the same event, mode and granularity is not specified

- it is a source of non determinism
- \sum When an error occurs
 - rollback of all operations performed by the triggers
 - rollback of the triggering statement in the triggering transaction



Non termination

 \sum Trigger execution may activate other triggers

- Cascaded trigger activation may lead to non termination of trigger execution
- \sum A maximum length for the cascading trigger execution may be set
 - default = 32 triggers
- ${} \boxdot$ If the maximum is exceeded
 - an execution error is returned



Mutating tables

▷ A *mutating table* is the table modified by the statement (i.e., event) triggering the trigger
 ▷ The mutating table

- *cannot* be accessed in row level triggers
- may *only* be accessed in statement triggers
- \sum Limited access on mutating tables only characterizes Oracle applications
 - accessing mutating tables is *always* allowed in SQL3



Example

 \sum Trigger to manage reorder in an inventory stock

- when the stocked quantity of a product goes below a given threshold
- a new order for the product should be issued

The following database schema is given Inventory (Part#, QtyOnHand, ThresholdQty, ReorderQty)
PendingOrders(Part#, OrderDate, OrderedQty)



Example

 \sum Trigger to manage reorder in an inventory stock

- when the stocked quantity of a product goes below a given threshold
- a new order for the product should be issued
- \supset Event
 - Update of the quantity on hand for product x
 - Insert of a new product x
- \supset Execution semantics
 - After the modification event
 - Separate execution for each row of the Inventory table



Trigger example

CREATE TRIGGER Reorder AFTER UPDATE OF QtyOnHand OR INSERT ON Inventory FOR EACH ROW



Example

 \sum Trigger to manage reorder in an inventory stock

- when the stocked quantity of a product goes below a given threshold
- a new order for the product should be issued
- \supset Condition
 - The quantity on hand is below a given threshold



Trigger example

CREATE TRIGGER Reorder AFTER UPDATE OF QtyOnHand OR INSERT ON Inventory FOR EACH ROW WHEN (NEW.QtyOnHand < NEW.ThresholdQty)



Example

 \sum Trigger to manage reorder in an inventory stock

- when the stocked quantity of a product goes below a given threshold
- a new order for the product should be issued
- \supset Condition
 - The quantity on hand is below a given threshold *and* there are no pending orders for product x
 - This part cannot be introduced into the WHEN clause
- \supset Action
 - Issue an order with given reorder quantity for product x

Example: Trigger body

```
DECLARE
N number;
BEGIN
select count(*) into N
from PendingOrders
where Part # = :NEW.Part #;
If (N=0) then
  insert into PendingOrders(Part#,OrderedQty,OrderDate)
  values (:NEW.Part#, :NEW.ReorderQty, SYSDATE);
end if;
END;
```



Complete trigger example

CREATE TRIGGER Reorder AFTER UPDATE OF QtyOnHand OR INSERT ON Inventory FOR EACH ROW WHEN (NEW.QtyOnHand < NEW. ThresholdQty) DECLARE N number; **BEGIN** select count(*) into N from PendingOrders where Part # = :NEW.Part #;If (N=0) then insert into PendingOrders(Part#,OrderedQty,OrderDate) values (:NEW.Part#, :NEW.ReorderQty, SYSDATE); end if; END;





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DB2 Triggers



Trigger syntax

CREATE TRIGGER TriggerName

Mode Event

ON TargetTable

[REFERENCING ReferenceName]

FOR EACH Level

WHEN Predicate Procedural SQL Statements

Mode is BEFORE or AFTER
 Event is INSERT or DELETE or UPDATE
 Only one event is allowed for a single trigger
 Level is ROW or STATEMENT

Trigger syntax

CREATE TRIGGER *TriggerName Mode Event* ON *TargetTable* [REFERENCING *ReferenceName*] FOR EACH *Level* WHEN *Predicate*

Procedural SQL Statements

The condition may be specified for *both* row and statement triggers



Trigger syntax

CREATE TRIGGER *TriggerName Mode Event* ON *TargetTable* [REFERENCING *ReferenceName*] FOR EACH *Level* WHEN *Predicate Procedural SQL Statements*

 \sum State variables are available for *both* row and statement triggers

- OLD and NEW tuple variables for row triggers
- OLD_TABLE and NEW_TABLE set variables for statement triggers



Trigger semantics

- \supset Execution modes
 - immediate before
 - immediate after
- Σ Granularity is
 - row (tuple)
 - statement
- \sum Execution is triggered by insert, delete, or update statements in a transaction



Trigger semantics

 \sum Before triggers cannot modify the database

- apart from the tuples affected by the triggering statement
 - tuple variables are used
- cannot trigger other triggers
- ${}^{\textstyle \sum}$ The execution of row and statement triggers with the same mode is in arbitrary order
- ${}^{\textstyle \sum}$ When more triggers are activated on the same event and mode
 - they are executed in *creation order*

 \sum Trigger execution is *deterministic*

Trigger semantics

- ${\ensuremath{\unrhd}}$ Cascading trigger execution is allowed up to a maximum number of triggers in the execution chain
- \supset When an error occurs
 - rollback of all operations performed by the triggers
 - rollback of the *entire transaction*



Execution algorithm

- \sum Transaction T contains a statement S which generates event E
- 1. T's execution is suspended and its state is saved into a stack
- 2. Old and new values of E are computed
- 3. Before triggers on E are executed
- 4. New values are applied to the DB (the modification due to E is actually performed)
 - Constraints are checked
 - compensative actions may trigger other triggers, hence cause a recursive invocation of the same execution procedure



Execution algorithm

- 5. After triggers triggered by E are executed
 - If any trigger contains an action A which triggers other triggers
 - the same execution procedure is recursively invoked on A
- 6. The execution state of T is extracted from the stack and T is resumed



\supset The trigger

- monitors the Inventory table
- Inserts into an audit table the information on
 - the user performing updates on the table
 - the update date and number of updated tuples
- \sum The following table is given

InventoryAudit (UserName, Date, Update#)



\supset Event

- Update of the Inventory table
- \supset Execution semantics
 - After the modification event
 - Separate execution for each update instruction
 - Statement semantics
- \supset No condition for execution



Trigger example

CREATE TRIGGER UpdateAudit AFTER UPDATE ON Inventory FOR EACH STATEMENT insert into InventoryAudit (UserName, Date, Update#) values (USER, SYSDATE, (select count(*) from OLD_TABLE));





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Comparing Oracle and DB2 Triggers



Differences between Oracle and DB2

		Oracle	DB2
	Reference to Old_Table and New_Table in statement triggers	No	Yes
	When clause in statement triggers	No	Yes
	Execution order between row and statement triggers with same mode	Specified	Arbitrary
	Execution order between triggers with same event, mode and granularity	Unspecified	Creation Order
	More than one triggering event allowed	Yes	No
	Forbidden access to the mutating table	Yes for row triggers	No
	Availability of the instead semantics	Yes	No
D <mark>B</mark> G	Database modifications allowed in before triggers	Yes	Only NEW variables



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Trigger Design



Trigger design

${\ensuremath{\unrhd}}$ The design of a single trigger is usually simple

• Identify

- execution semantics
- event
- condition (optional)
- action



Trigger design

 \sum Understanding *mutual* interactions among triggers is more complex

- The action of one trigger may be the event of a different trigger
 - Cascaded execution
- \sum If mutual triggering occurs
 - Infinite execution is possible



Trigger execution properties

\supset Termination

 For an arbitrary database state and user transaction, trigger execution *terminates* in a final state (also after an abort)

\sum Confluence

 For an arbitrary database state and user transaction, trigger execution *terminates in a unique final state*, independently of the execution order of triggers

 $\mathop{\textstyle\sum}$ Termination is the most important property

 $\mathop{\textstyle \sum}$ Confluence is enforced by deterministic trigger execution



Guaranteeing termination

- \sum Termination is guaranteed at run time by aborting trigger execution after a given cascading length
- ${\ensuremath{\unrhd}}$ Termination may be verified at design time by means of the triggering graph
 - a node for each trigger
 - a directed edge $T_i \rightarrow T_j$ if trigger T_i is performing an action triggering trigger T_j
- \sum A cycle in the graph shows potential non terminating executions



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\sum Trigger managing salary amounts

- When a given average salary value is exceeded, a salary reduction is automatically enforced
- The following table is given Employee (<u>Emp#</u>, Ename, ..., Salary)
- ∑ Event
 - Update of the Salary attribute in Employee
 - Insert into Employee
 - Will write only trigger for update



\sum Trigger managing salary amounts

- When a given average salary value is exceeded, a salary reduction is automatically enforced
- The following table is given Employee (Emp#, Ename, ..., Salary)
- \supset Execution semantics
 - After the modification events
 - Separate execution for each update instruction
- \sum No condition for execution



CREATE TRIGGER SalaryMonitor AFTER UPDATE OF Salary ON Employee FOR EACH STATEMENT BEGIN update Employee set Salary = Salary * K where 2500 < (select AVG (Salary) from Employee); END;

The value of K may be

- $K = 0.9 \rightarrow$ execution terminates
- $K = 1.1 \rightarrow \text{infinite execution}$

SalaryMonitor



Trigger applications

\sum Internal applications

- maintenance of complex integrity constraints
- replication management
- materialized view maintenance
- \supset Business Rules
 - Incorporate into the DBMS application knowledge
 - E.g., reorder rule

 \supset Alerters

• widely used for notification



Triggers for constraint management

- ${\ensuremath{\unrhd}}$ Triggers are exploited to enforce complex integrity constraints
- \supset Design procedure
 - 1. Write the constraint as a SQL predicate
 - It provides a condition for the trigger execution
 - 2. Identify the events which may violate the constraint
 - i.e. the condition
 - 3. Define the constraint management technique in the action



\sum The following tables are given

- Supplier S (<u>S#</u>, SName, ...)
- Part P (<u>P#</u>, PName, ...)
- Supply SP (<u>S#</u>, <u>P#</u>, Qty)
- \supset Constraint to be enforced
 - A part may be supplied by at most 10 different suppliers



\sum Constraint predicate

select P# from SP group by P# having count(*) > 10

set of parts violating the constraint

 \sum Events

- insert on SP
- update of P# on SP
- \sum Action
 - reject the violating transaction

\supset Execution semantics

- *after* the modification
- statement level
 - to capture the effect of the entire modification
 - (Oracle) to allow access to the mutating table

\sum (Oracle) No condition

- The condition cannot be specified in the WHEN clause
- It is checked in the trigger body
- \sum Design for Oracle trigger semantics



```
CREATE TRIGGER TooManySuppliers
AFTER UPDATE OF P# OR INSERT ON SP
DECLARE
N number;
BEGIN
select count(*) into N
from SP
where P# IN (select P# from SP
              group by P#
              having count(^{*}) > 10);
if (N <> 0) then
  raise_application_error (xxx, 'constraint violated');
end if;
END;
```

$\mathop{\textstyle\sum}$ The following tables are given

- Supplier S (<u>S#</u>, SName, ...)
- Part P (<u>P#</u>, PName, ...)
- Supply SP (<u>S#</u>, <u>P#</u>, Qty)

 \sum Constraint to be enforced

- The quantity of a product supply cannot be larger than 1000. If it is larger, trim it to 1000.
- \sum Check constraints do not allow compensating actions
 - Implement with a trigger



\sum Constraint predicate

- Qty > 1000
- It is also the trigger condition
- \sum Events
 - insert on SP
 - update of Qty on SP
- \sum Action
 - Qty = 1000



\sum Execution semantics

- *before* the modification takes place
 - its effect can be changed before the constraint is checked
- row level
 - each tuple is modified separately



CREATE TRIGGER ExcessiveQty BEFORE UPDATE OF Qty OR INSERT ON SP FOR EACH ROW WHEN (NEW.Qty > 1000) BEGIN :NEW.Qty := 1000; END;



Triggers for materialized view maintenance

 ${\ensuremath{\unrhd}}$ Materialized views are queries persistently stored in the database

- provide increased performance
- contain redundant information
 - e.g., aggregate computations

 \sum Triggers are exploited to maintain redundant data

 Propagate data modifications on tables to materialized view



Σ Tables

- Student S (<u>SId</u>, SName, DCId)
- Degree course DC (<u>DCId</u>, DCName)
- \sum Materialized view
 - Enrolled students ES (<u>DCId</u>, TotalStudents)
 - For each degree course, TotalStudents counts the total number of enrolled students
 - Defined by query

SELECT DCId, COUNT(*) FROM S GROUP BY DCId;



Design example (3)

Σ Tables

- Student S (<u>SId</u>, SName, DCId)
- Degree course DC (<u>DCId</u>, DCName)
- \supset Materialized view
 - Enrolled students ES (<u>DCId</u>, TotalStudents)
 - For each degree course, TotalStudents counts the total number of enrolled students
 - A new degree course is inserted in materialized view ES when the first student is enrolled in it
 - A degree course is deleted from ES when the last student quits it



Design example (3)

\square Database schema

S (<u>SId</u>, SName, DCId) DC (<u>DCId</u>, DCName) ES (<u>DCId</u>, TotalStudents)

- \sum Propagate modifications on table S to materialized view (table) ES
 - Inserting new tuples into S
 - Deleting tuples from S
 - Updating the DCId attribute in one or more tuples of S



Design example (3)

 \sum Design three triggers to manage separately each data modification

- Insert trigger, delete trigger, update trigger
- All triggers share the same execution semantics
- \supset Execution semantics
 - *after* the modification takes place
 - Table ES is updated after table S has been modified
 - row level
 - Separate execution for each tuple of table S
 - significantly simpler to implement



Insert trigger (3)

\supset Event

- insert on S
- \sum No condition
 - It is always executed
- \supset Action
 - if table ES contains the DCId in which the student is enrolled
 - increment TotalStudents
 - otherwise
 - add a new tuple in table ES for the degree course, with TotalStudents set to 1



Insert trigger (3)

CREATE TRIGGER InsertNewStudent AFTER INSERT ON S FOR EACH ROW DECLARE N number; BEGIN --- check if table ES contains the tuple for the degree --- course NEW.DCId in which the student enrolls select count(*) into N from ES where DCId = :NEW. DCId;



Insert trigger (3)

```
if (N <> 0) then
    --- the tuple for the NEW.DCId degree course is
    --- available in ES
    update ES
    set TotalStudents = TotalStudents +1
    where DCId = :NEW.DCId;
else
    --- no tuple for the NEW.DCId degree course is
    --- available in ES
    insert into ES (DCId, TotalStudents)
    values (:NEW.DCId, 1);
end if;
END;
```



Delete trigger (3)

\supset Event

- delete from S
- \sum No condition
 - It is always executed
- \sum Action
 - if the student was the only student enrolled in the degree course
 - delete the corresponding tuple from ES
 - otherwise
 - decrement TotalStudents



Delete trigger (3)

CREATE TRIGGER DeleteStudent AFTER DELETE ON S FOR EACH ROW DECLARE N number; BEGIN --- read the number of students enrolled on --- the degree course OLD.DCId select TotalStudents into N from ES where DCId = :OLD.DCId;



Delete trigger (3)

```
if (N > 1) then
    --- there are many enrolled students
    update ES
    set TotalStudents = TotalStudents -1
    where DCId = :OLD.DCId;
else
    --- there is a single enrolled student
    delete from ES
    where DCId = :OLD.DCId;
end if;
END;
```



\supset Event

- Update of DCId on S
- \sum No condition
 - It is always executed
- \supset Action
 - update table ES for the degree course where the student was enrolled
 - decrement TotalStudents, or delete tuple if last student
 - update table ES for the degree course where the student *is currently* enrolled
 - increment TotalStudents, or insert new tuple if first student



CREATE TRIGGER UpdateDegreeCourse AFTER UPDATE OF DCId ON S FOR EACH ROW DECLARE N number; **BEGIN** --- read the number of students enrolled in --- degree course OLD.DCId select TotalStudents into N from ES where DCId = :**OLD**.DCId;



```
if (N > 1) then
    --- there are many enrolled students
    update ES
    set TotalStudents = TotalStudents -1
    where DCId = :OLD.DCId;
else
    --- there is a single enrolled student
    delete from ES
    where DCId = :OLD.DCId;
end if;
```



--- check if table ES contains the tuple for the degree --- course NEW.DCId in which the student is enrolled select count(*) into N from ES where DCId = :NEW. DCId;



```
if (N <> 0) then
 --- the tuple for the NEW.DCId degree course is available in ES
       update ES
       set TotalStudents = TotalStudents +1
       where DCId = :NEW.DCId;
else
 --- no tuple for the NEW.DCId degree course is available in ES
      insert into ES (DCId, TotalStudents)
      values (:NEW.DCId, 1);
end if;
```

```
D<mark>B</mark>G
```

END;