

Introduction to databases

Database design

Database design

- \sum Entity-Relationship model
- \sum Conceptual design
- Σ Logic design
- \sum Normalization



Database design

Entity-Relationship Model

Entity-Relationship model

- \sum Life cycle of an information system
- Σ Databases design
- Σ Entities and relationships
- \supset Attributes
- \supset Identifiers
- \supset Generalization
- \sum E-R schema documentation
- \supset UML and E-R



Entity-Relationship model

Databases design

 \sum The design of a database is one of the activities in the development of an information system

 must be considered in the broader context of the life cycle of an information system

Feasibility study

\sum Feasibility study

 determination of the costs of the various alternatives and of the implementation priorities of the system components



 \sum Collection and analysis of the requirements

- definition of properties and functions of the information system
- requires interaction with the user
- produces a complete but informal description of the system to be implemented



\sum Design

- divided into data and application design
- produces formal descriptions



\supset Implementation

• creation of the information system according to the characteristics defined in the design phase



\sum Validation and testing

- verification of the correct functioning and quality of the information system
- can lead to changes in requirements or a revision of the design



\sum Functioning

- system operation
- requires management and maintenance operations



\sum Prototyping

- rapid creation of a simplified version of the system in order to evaluate its characteristics
- can lead to changes in requirements or a revision of the design







Entity-Relationship model

Database design

Database design

- ${\ensuremath{\unrhd}}$ The database is an important component of the overall system
- \sum Data-driven design methodology
 - database design precedes the design of the applications that use it
 - greater attention to the design phase than the other phases

Design methodology

\sum A design methodology consists of

- decomposition of the project activity into successive independent steps
- strategies to be followed in the various steps and criteria for choosing strategies
- reference models to describe the input and output data of the various phases

Design methodology: Example

\supset Athletic training

Activity decomposition

 physical condition
 enhancement
 velocity

Design methodology: Example

\supset Athletic training

- activity decomposition
- strategies to follow in the various steps
 - 1. A) diet
 - B) exercises to reduce the percentage of body fat
 - 2a.A) strength exercises
 - B) resistance exercises

Design methodology: Example

\supset Athletic training

- activity decomposition
- strategies to follow in the various steps
- reference models to describe the input and output data of the various phases
 - input data: current weight, % of body fat output data: model of the body structure of a fit person
 - 2a. input data: fit person model
 - output data: body structure model of the average athlete

Properties of the methodology

\supset Generality

- same methodology regardless of the problem and the tools available
- \sum Quality of the result
 - in terms of correctness, completeness and efficiency with respect to the resources used
- \sum Ease of use
 - of both strategies and reference models

Data-driven design

 ${\ensuremath{\unrhd}}$ For databases, methodology based on the separation of decisions

- what to represent in the database
 - conceptual design
- *how* to represent it
 - Iogical and physical design

Stages of database design





Application requirements

- Informal specifications of the reality of interest
 - application properties
 - application functionalities

Stages of database design





Conceptual design

 \sum Representation of informal specifications in the form of a *conceptual schema*

- formal and complete description, which refers to a conceptual model
- independence from implementation aspects (data model)
- the target is the representation of the *information content* of the database

Stages of database design



Logical design

- \sum Translation of the conceptual schema into the logical schema
 - refers to the chosen logical data model
 - criteria are used to optimize the operations which must be performed on the data
 - quality of the schema verified by formal techniques (normalization)
Stages of database design



Physical design

 \sum Specification of physical data storage parameters (organization of files and indexes)

produces a physical model, which depends on the chosen DBMS

Stages of database design





Entity-Relationship model

Entities and relationships

E-R model (Entity-Relationship)

- ${\ensuremath{\unrhd}}$ It is the most widespread conceptual model
- \sum Provides constructs to describe data structure specifications
 - in a simple and understandable way
 - with graphic formalism
 - regardless of the data model, which can be chosen later
- \sum There are numerous variations

Main constructs of the E-R model

- \supset Entities
- Σ Relationships
- \supset Attributes
- \supset Identifiers
- $\mathop{\textstyle \textstyle \sum}$ Generalizations and subsets



- \sum Entities represent classes of real-world objects (people, things, events, ...), which have
 - common properties
 - autonomous existence
- \sum Examples: employee, student, item
- \sum An occurrence of an entity is an object of the class that the entity represents

Relationship

Relationship name

- \sum Represents a logical link between two or more entities
- \sum Examples: exam between student and course, residence between person and municipality
- $\mathop{\textstyle \sum}$ Not to be confused with the relationship of the relational model
 - sometimes it is named association

Student







Person

Municipality





Occurrences of a relationship







Occurrences of a relationship

An occurrence of a relationship is an n-tuple (pair in the case of a binary relationship) consisting of occurrences of entities, one for each of the entities involved



Occurrences of a relationship

 \sum An occurrence of a relationship is an n-tuple (pair in the case of a binary relationship) consisting of occurrences of entities, one for each of the entities involved

 \sum No identical n-tuples are allowed



Student

- \sum Cardinalities are specified for each entity participating in a relationship
- \sum They describe the maximum and minimum number of relationship occurrences to which an entity occurrence can participate
 - minimum cardinality
 - 0 (optional participation)
 - I (mandatory participation)
 - Maximum cardinality
 - 1 (at most one occurrence)
 - N (arbitrary number of occurrences)

> 1 to 1 correspondence



Professor



















 \sum N to N correspondence



Student









Limitations of a binary relationship



${}^{\textstyle \sum}$ It is not possible that a student takes the same exam more than once

Ternary relationship

 \sum A student can take the same exam more than once at different times

 \sum Example of an exam instance

$$\begin{array}{ccc} s_1 & c_1 & t_1 \\ s_1 & c_1 & t_2 \end{array}$$

Ternary relationship



 \sum A student can take the same exam more than once at different times

 \sum Example of an exam instance










Cardinality of ternary relationships



Observations

- \sum Minimum cardinalities are rarely 1 for all entities involved in a relationship
- \sum The maximum cardinalities of an n-ary relationship are (practically) always N
 - if the participation of an entity E has a maximum cardinality of 1, it is possible to eliminate the n-ary relationship and link entity E with the others through binary relationships







- \sum Relationship between an entity and itself
- \sum If the relationship is not symmetrical, the two roles of the entity must be defined



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\sum An employee might have several superiors



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Entity-Relationship model

Attributes

The attribute

Name of the attribute

- \sum It describes an elementary property of an entity or a relationship.
- \sum Examples
 - Surname, name, student id are attributes that describe the entity student.
 - Grade is an attribute that describes the relationship exam.
- \sum Each attribute is characterized by the *domain*, the set of eligible values for the attribute.



















Composite attribute



 \sum Group of attributes that have closely connected meanings or uses.

Composite attribute



- \sum Group of attributes that have closely connected meanings or uses.
- \sum Example



- ${\ensuremath{\unrhd}}$ It can be specified for the attributes of entities or relationships
- \sum It describes the minimum and maximum number of attribute's values associated to an instance of an entity or a relationship.
 - If omitted, it corresponds to (1,1)
 - minimum 0 corresponds to having an attribute that admits the null value
 - maximum N corresponds to having an attribute that can take more than one value for the same occurrence (multivalued attribute)













The Entity-Relationship Model

Identifiers

Identifier

\sum It is specified for each entity

- \sum It describes concepts (attributes and/or entities) of the scheme that allow to identify uniquely the instances of an entity.
 - Each entity must have at least one identifier
 - There can be be more than one appropriated identifier for a given entity.

Internal Identifier

\sum Simple: consisting of only one attribute



Internal Identifier

\sum Simple: consisting of only one attribute



 \sum Composite: consisting of multiple attributes












 \supset One entity that does not have sufficient internal attributes able to define an identifier is called *weak entity*.



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 \sum A weak entity must partecipate with cardinality (1,1) in each of the relationships that provide part of the identifier







 \sum Is it possible to represent in the same order more order lines for the same product?



Remarks

 $\mathop{\textstyle \sum}$ An external identifier may involve an entity that is itself externally identified

Identification cycles must not be generated



Remarks

Relationships do not have identifiers Student ID Course code Name Surname O Student (0,N) Course Name O Student (0,N) Course



The Entity-Relationship Model

Generalizations

Generalization



- Σ It describes a logical link between an entity E and one or more entities E₁, E₂,..., E_n, that are particular cases of E.
 - E is called parent entity, and is a generalization of E₁, E₂,..., E_n
 - E₁, E₂,..., E_n are called child entities, and are specialization of E













Generalization: property

- \sum Each instance of a child entity is also an instance of the parent entity.
- \sum Each property of the parent entity (attributes, identifiers, relationships, other generalizations) is also a property of each child entity.
 - Property known as *inheritance*
- \sum One entity can be involved in several different generalizations.











Generalization: property

\supset Orthogonal characteristics

- *total* generalization if each instance of the parent entity is an instance of at least one of the child entities, *partial* otherwise.
- *exclusive* if each instance of the parent entity is at most one instance of one of the child entities, *overlapping* otherwise.









Subset

- \sum Particular case of generalization with only one child entity
 - the generalization is always partial and exclusive.





The Entity-Relationship Model

Documentation of E-R schemes

Documentation of E-R schemes



Documentation of E-R schemes

\sum Data Dictionary

 allows to enrich the E-R scheme with natural language description of entities, relationships and attributes

Data dictionary: example

Entity	Description	Attributes	Identifier
Student	University student	Student ID, Surname, Name, CFU acquired, Grades average	Student ID
Professor	University professor	Professor ID, Department, Surname, Name	Professor ID
Course	Courses offered by the university	Course code, Name, CFU	Course code
Time	Dates on which exams were taken	Date	Date 143

Data dictionary: example

Relationship	Description	Entities involved	Attributes
Exam	It associates a student to the exams taken and memorizes the mark obtained	Student(0,N),Course(0,N),Time(1,N)	Grade
Holder	It associates each course to the professor who teachers the course.	Course (1,1), Professor (0,N)	

Documentation of E-R schemes

\sum Data Dictionary

 allows to enrich the E-R scheme with natural language description of entities, relationships and attributes

\sum Integrity constraints on data

- may not always be explicitly indicated in an E-R scheme
- can be described in natural language

Constraints of integrity on data: example

Constraints of integrity		
RV1	The grade of an exam can only take values between 0 and 30	
RV2	Each student cannot pass the same exam twice	
RV3	A student may not take more than three exams for the same course during the same academic year	

Documentation of E-R schemes

\sum Data Dictionary

 allows to enrich the E-R scheme with natural language description of entities, relationships and attributes

\sum Data integrity constraints

- may not always be explicitly indicated in an E-R scheme
- can be described in natural language
- \sum Rules for deriving data
 - allow to define how a scheme concept can be obtained (by logical inference or arithmetic calculation) from other scheme concepts

Rules for deriving data: example

Derivation rules			
RD1	The number of credits acquired by a student is obtained by adding the number of credits of the courses for which the student has passed the exam		
RD2	The average mark is obtained by calculating the average marks of the exams passed by a student		



The Entity-Relationship Model

UML and E-R

UML and E-R

>> UML (Unified Modeling Language)

- is a model of a software application
 - structural and behavioural aspects (data, operations, processes and architectures)
- rich formalism
 - Diagrams of classes, of actors, of sequence, of communication, of states,...
- Σ E-R
 - is a model of a database
 - Structural aspects of an application
 - specific formalism for database modelling

UML and E-R

 \sum Main charachteristics of UML that differs with respect to E-R diagrams

- absence of standard notation to define identifiers
- possibility to add notes to comment on diagrams
- possibility to indicate the navigation direction of an association (not relevant in the design of a database)

UML and E-R

\supset Different formalisms

- \sum The class diagram of an application is different from the E-R scheme of the database
- \sum The class diagram, even if designed for a different use, may be adapted for the description of the conceptual design of a database